

## 642:573:01 MISCELLANEOUS EXERCISES (ASSIGNMENT #4)

**1.** Let  $f$  be a function defined on  $[a, b]$  and  $a = x_0 < \cdots < x_n = b$  be nodes on the same interval. In each subinterval  $[x_{i-1}, x_i]$  introduce two new nodes  $x_{i-2/3}, x_{i-1/3}$  which are  $1/3, 2/3$  of the way from left to right respectively. The **Bessel interpolant** of  $f(x)$  in this scheme is the piecewise-cubic-polynomial function which is a cubic polynomial on each  $[x_{i-1}, x_i]$  and interpolates  $f(x)$  on that interval at the four points  $x_{i-1}, x_{i-2/3}, x_{i-1/3}$  and  $x_i$ . Denote the Bessel interpolant of  $f(x)$  on  $[a, b]$  by  $C(x)$ .

**A.** Assuming  $f(x)$  has four derivatives, recall that if  $x \in [x_{i-1}, x_i]$  then

$$f(x) - C(x) = \frac{1}{4!} f^{(4)}(\xi_i)(x - x_{i-1})(x - x_{i-2/3})(x - x_{i-1/3})(x - x_i)$$

where  $x_{i-1} < \xi_i < x_i, i = 1, \dots, n$ . Deduce that if the nodes are equally spaced with  $x_i - x_{i-1} = h$ , then

$$|f(x) - C(x)| \leq \frac{M_4}{4!} \cdot \frac{h^4}{81},$$

where  $M_4 = \max\{|f^{(4)}(x)| : a \leq x \leq b\}$ , is a bound for the error of approximation that is valid for all  $x \in [a, b]$ .

**B.** Apply the result of **A** above to determine what choice of  $h$  will be required in order that the Bessel interpolant of the function

$$f(x) = 24e^{3(x-1)}$$

on  $[-1, 1]$  approximate it “uniformly on  $[-1, 1]$  with error  $< 10^{-8}$ ,” *i.e.*, such that for every  $x \in [-1, 1]$  the error is bounded by 0.00000 001. Compare this with the number of interpolation points required if the zeros of a Chebyshev polynomial  $T_n(x)$  are used to interpolate  $f(x)$  in the usual way.

**C.** Are any of the derivatives of  $C(x)$  continuous, in general?

**2.** Consider the single-step Method (the **simple Kutta method**) given by

$$F(x, y, h; f) = \frac{1}{6} [k_1 + 4k_2 + k_3]$$

$$\text{where } k_1 = f(x, y)$$

$$k_2 = f\left(x + \frac{h}{2}, y + \frac{h}{2} k_1\right)$$

$$k_3 = f(x + h, y + h \cdot (2k_2 - k_1)).$$

Note the obvious relation with Simpson’s rule of approximate quadrature (more apparent here than in the Runge-Kutta method). **Show** that this is a third-order Method in the sense that

$$\tau(x, Y, h; f) = O(h^3)$$

when  $Y(x)$  is a solution of  $Y'(x) = f(x, Y(x))$ . (Use a symbolic manipulation program—don’t make your life harder than it has to be . . . .) Comment on the stability of this Method.

**3.** Let  $[a, b]$  be a real interval. Suppose we have a function  $E$  whose domain is the set of closed subintervals of  $[a, b]$ —so this is a “set function”—and which takes real values. Suppose  $e(h)$  is a continuous nonnegative-real-valued function of the positive real variable  $h$  which “estimates  $E$ ” in the sense that

$$|E([c, d])| \leq e(d - c)$$

for each interval  $[c, d] \subseteq [a, b]$ . Suppose the estimation function satisfies

$$\lim_{h \rightarrow 0^+} \frac{e(h)}{h} = 0 \quad \text{and} \quad \lim_{h \rightarrow +\infty} \frac{e(h)}{h} = \infty.$$

(a) Let  $\epsilon > 0$  be given, and produce a sequence of points

$$a = x_0 < x_1 < \dots$$

in the following way: if  $x_k$  has been constructed, then  $x_{k+1}$  is chosen as  $x_k + h_{k+1}$  where  $h_{k+1}$  is the largest solution  $h_{k+1} > 0$  of

$$e(h_{k+1}) = \frac{\epsilon}{b-a} h_{k+1}$$

having the property that if  $h < h_{k+1}$  then also

$$e(h) \leq \frac{\epsilon}{b-a} h$$

unless such an  $x_{k+1}$  would be  $> b$ , in which case one chooses  $x_{k+1} = b$ . **Show**, using the standard properties of the real numbers, that such an  $h_{k+1}$  always exists.

(b) **Show** that the process described in (a) always terminates; *i.e.*, at some step *it must happen that*  $x_{k+1} = b$ . With  $n$  denoting the index for which  $x_n = b$ , **show** that then

$$\sum_{k=0}^{n-1} |E[x_k, x_{k+1}]| \leq \epsilon.$$

c) The result of (a) and (b) can be used to provide theoretical justification for adaptive methods of quadrature (for integrals) and adaptive methods for step-size control in solving initial-value problems of o. d. e. For example, in adaptive approximate integration of  $f(x)$  over  $[a, b]$  using a quadrature method which assigns  $I[c, d]$  as the approximate value of the integral  $\int_c^d f(x) dx$  with an error  $E[c, d]$ —*e.g.*, for the

midpoint method  $E[c, d] = (d - c)^3 \frac{f''(\xi)}{24}$ —for which an estimate of the type considered in (a) and (b) is available, one proceeds as follows. Decide to accept a total error of  $\epsilon$ . Pick  $x_1$  so that applying the method on  $[a, x_1]$  incurs an  $|\text{error}| \leq h_1 \cdot \epsilon / (b - a)$ ; then pick  $x_2$  so that applying the method on  $[x_1, x_2]$  incurs an  $|\text{error}| \leq h_2 \cdot \epsilon / (b - a), \dots$  **Explain** how the result of (a) and (b) guarantees that one will eventually reach  $x_n = b$  in finitely many steps, thus theoretically excluding the embarrassing possibility that the adaptive quadrature process will loop forever without ever reaching  $b$ . {A similar application, though it's a bit more complicated, applies to adaptive step-size choice in approximate solution of o. d. e. initial-value problems as in the o. d. e. notes.}