# Random Polygon 

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An initial "random polygon" on the plane with $n$ vertices is equivalent to the data of a vector $v \in \mathbb{C}^{n}$. The averaging process is equivalent to applying the matrix

$$
A=\frac{1}{2}(I+T)
$$

on $v$ where $I$ is the identity matrix and $T$ is the standard shift matrix given as $T_{i j}=\delta_{i, j+1}$ where the indices are interpreted mod $n$. Therefore averaging $N$ times simply amounts to taking the entries of $A^{N} v$ and plotting them on the complex plane.

A crucial role is played by the eigenvalues and eigenvectors of $A$. The diagonalization of $A$ simply amounts to diagonalizing $T$ which is well-known: One notes the matrix equation $T^{n}=1$ which is actually the characteristic polynomial of $T$. Thus the eigenvalues satisfy $\omega^{n}=1$, that is, they're simply the $n$th roots of unity. The corresponding eigenvectors are also easy to write down. Let $\omega=e^{2 \pi i / n}$. Then the (normalized) eigenvector corresponding to $\omega^{k}$ ( $k=0, \ldots, n-1$ ) is simply

$$
\left(e_{k}\right)_{j}=\frac{1}{\sqrt{n}} \omega^{k j} .
$$

$A$ will have the same eigenvectors $e_{k}$ with corresponding eigenvalues given by

$$
\lambda_{k}=\frac{1}{2}\left(1+\omega^{k}\right)=e^{i \pi k / n} \cos \left(\frac{\pi k}{n}\right) .
$$

From this one can easily see that the largest eigenvalue is given by $k=0$ when $\lambda_{0}=1$ and the eigenvector is simply given by $\left(e_{0}\right)_{j}=\frac{1}{\sqrt{n}}$. The others have magnitude strictly less than one with the second largest one corresponding to $k= \pm 1$. Now $T$ being a unitary matrix means that the eigenvectors form an orthogonal basis for $\mathbb{C}^{n}$ and so we can easily expand $v$ in terms of this eigenbasis

$$
v=\sum_{j=0}^{n-1}\left\langle e_{j}, v\right\rangle e_{j}
$$

where we note that the $j$-th coefficient given by the inner product $\left\langle e_{j}, v\right\rangle$ is simply given by the discrete Fourier transform

$$
x_{j}:=\left\langle e_{j}, v\right\rangle=\frac{1}{\sqrt{n}} \sum_{k=1}^{n} v_{j} e^{-\frac{2 \pi i k j}{n}}
$$

The action of $A^{N}$ now reads

$$
A^{N} v=\sum_{j=0}^{n-1} x_{j}\left(\lambda_{j}\right)^{N} e_{j}:=\sum_{j=0}^{n-1}\left(\lambda_{j}\right)^{N} \alpha_{j}
$$

Let us now interpret each of the terms in decreasing absolute value of $\lambda_{j}^{N}$. First, $\alpha_{0}$ is simply the vector with all entries $\frac{1}{n}\left(v_{1}+\ldots+v_{n}\right)$ which corresponds to the center of the initial random polygon. We can freely choose this to be zero. The subleading term i.e the term

$$
E:=\left(\lambda_{1}\right)^{N} \alpha_{1}+\left(\lambda_{n-1}\right)^{N} \alpha_{n-1}
$$

is what will play the crucial role. Writing it out explicitly we have

$$
E_{m}=\frac{1}{\sqrt{n}}\left(\cos \left(\frac{\pi}{n}\right)\right)^{N}\left(x_{1} e^{\frac{i \pi N}{n}} e^{\frac{2 \pi i m}{n}}+x_{n-1} e^{\frac{-i \pi N}{n}} e^{\frac{-2 \pi i m}{n}}\right)
$$

where $x_{1}, x_{n-1}$ are components of the discrete Fourier transform as given above. The rest of the terms can be neglected for large $N$ and so $E$ is what should give us the sought-after ellipse. Indeed plotting this as in the program $E(L, n, N)$ this traces out precisely the ellipse matching the program IterPolygon1. It should be possible to figure out the properties of the ellipse such as its eccentricity and orientation from the coefficients $x_{1}$ and $x_{n-1}$, which are the only remnants of the initial vector $v$, although I haven't had the time to do that.

