Implementing and Experimenting with the Calabi-Wilf algorithm for random selection of a subspace over a finite field

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In fond memory of Eugenio Calabi (May 11, 1923 - September 25, 2023), and Herbert Saul Wilf (June 13, 1931 - January 7, 2012)

Abstract: In their beautiful note, "On the Sequential and Random Selection of Subspaces over a Finite Field", geometrical giant Eugenio Calabi and combinatorial giant Herbert Wilf proposed an elegant algorithm to do what is promised in their title. In the present note, written forty six years later, we describe a Maple package, written by the second author, that implements their beautiful algorithm, and then go on to report numerous experiments, performed by the first author, that demonstrate the efficiency and reliability of the Calabi-Wilf algorithm.

Memory Lane (by DZ)

One of my numerous temporary positions was during the academic year 1982-1983. Herbert Wilf, who later became my close collaborator, got me a one-year visiting position at the University of Pennsylvania. I occupied the office of David Harbater, who was on sabbatical that year, and it so happened that it was close to **both** Wilf's and Calabi's offices. I was in awe of both of them, since they were already legendary way back then. One day Herb mentioned to me that, while Calabi's expertise was almost *diametrically opposite* to his, they did have a *non-empty intersection*, and coauthored a cute little note [CW] that applies the *Wilf methodology* [W] for random selection of *combinatorial* objects, to something that besides being combinatorial (after all 'finite fields' are finite!) was also a *geometrical* object, that Calabi can relate to: namely **subspaces** of $GF(q)^n$.

Reminder about the Wilf approach for Random Generation of Combinatorial objects

The binomial coefficient,

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

famously counts the number of *subsets* with k elements of a 'universal' set of n elements, $\{1, \ldots, n\}$. It famously satisfies the *Pascal-Chu* recurrence

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

that breaks up all k-element subsets of $\{1, \ldots, n\}$ into those that **do** contain n, whose number is the first term on the right, and those that **do not** contain n, whose number is the second term on the right.

It follows that the *fraction* of k-element subsets of $\{1, \ldots, n\}$ that **do** contain n is

$$\frac{\binom{n-1}{k-1}}{\binom{n}{k}} = \frac{(n-1)!/((k-1)!(n-k)!)}{(n!/k!(n-k)!)} = \frac{k}{n}$$

,

hence the *probability* that a random k-subset of $\{1, \ldots, n\}$ does contain n is $\frac{k}{n}$, while the probability that it does not is $1 - \frac{k}{n}$.

This simple observation lead Wilf [W] to suggest the following *recursive* algorithm for the *random* generation, **uniformly at random**, of a k-subset of $\{1, \ldots, n\}$.

RandomSubset(n,k): if k < 0 or k > n return FAIL. If k = 0 then RETURN the empty set. Otherwise toss a loaded coin whose probability of Heads is k/n. If it lends on Heads, output RandomSubset(n-1,k-1) \cup {n}, else output RandomSubset(n-1,k).

This *philosophy* was used on many other combinatorial objects that posses *recursive structures*, and nicely described, and *implemented* (with FORTRAN source-code included!), in the Nijenhuis-Wilf classic [NW]. However, the Calabi-Wilf algorithm did not make it to that book. It was exposited in more detail by Igor Pak, in his insightful 'big picture' paper [P], written in his inimitable engaging style.

The Calabi-Wilf algorithm

The *q*-binomial coefficient,

$$\begin{bmatrix} n\\k \end{bmatrix} = \frac{[n]!}{[k]![n-k]!}$$

(where $[n]! := (1-q) \cdots (1-q^n)/(1-q)^n$ is the q-analog of n!), not-quite-as-famously, counts the number of k-dimensional subspaces of $GF(q)^n$. Each such subspace has many possible bases, of course, but only one in row-echelon form, so in order to count k-dimensional subspaces one should count $k \times n$ matrices over GF(q) in row-echelon form.

The right side of the *q*-Pascal-Chu recurrence

$$\begin{bmatrix} n \\ k \end{bmatrix} = \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} + q^k \begin{bmatrix} n-1 \\ k \end{bmatrix} ,$$

may be thought of (per [CW]) as counting complementary subsets of the totality of $k \times n$ matrices in echelon form:

• The first term counts those $k \times n$ matrices, $B = (b_{ij})_{1 \le i \le k, 1 \le j \le n}$ for which $b_{11} = 1$, $b_{i1} = 0$ $(1 < i \le k)$, $b_{1j} = 0$ $(1 < j \le n)$, and for which *necessarily* the remaining $(k-1) \times (n-1)$ array is a $(k-1) \times (n-1)$ matrix in row-echelon form.

• The second term counts the other basis matrices B, i.e. those whose first column is an **arbitrary** vector of length k, and for which the remaining $k \times (n-1)$ array is a basis matrix for a k dimensional subspace of $GF(q)^{n-1}$.

It follows that the *fraction* of matrices of the first kind is

$$\begin{bmatrix} n-1\\k-1 \end{bmatrix} / \begin{bmatrix} n\\k \end{bmatrix} = \frac{q^k - 1}{q^n - 1}$$

and it follows, in turn, that the *probability* that a random $k \times n$ matrix in row-echelon form is of the first kind is $\frac{q^k-1}{q^n-1}$.

This observation lead Calabi and Wilf [CW] to suggest the following *recursive* algorithm for the random generation, uniformly at random, of a $k \times n$ matrix over GF(q) in row-echelon form.

RandomSubspace(n,k): if k < 0 or k > n return FAIL. If k = 0 then RETURN the empty matrix (corresponding to the null subspace). Otherwise toss a loaded coin whose probability of Heads is $(q^k - 1)/(q^n - 1)$. If it lends on Heads, generate RandomSubspace(n-1,k-1), and convert it to a $k \times n$ matrix by putting on top the row vector 10^{n-1} , and prefixing the column vector $(10^{k-1})^T$ (where of course the sole 1 is shared by both).

Else generate RandomSubspace(n-1,k). Also generate a random column vector in GF(q), of length k, and prefix it at the front.

Maple Implementation

The Calabi-Wilf algorithm (and much more!) is implemented in the Maple package, CalabiWilf.txt, accompanying the present paper:

https://sites.math.rutgers.edu/~zeilberg/tokhniot/CalabiWilf.txt

It is procedure Rqnk(q,n,k). For example, typing

Rqnk(7,10,5);

would generate, uniformly-at-random, one of the

$$((7^{10}-1)(7^9-1)\cdots(7^6-1))/((7^5-1)(7^4-1)\cdots(7-1)) = 1602592475815614015216$$

 5×10 row-echelon matrices over GF(7). Here is one of them:

/5	6	3	2	5	1	0	0	0	0 \
6	1	0	0	6	0	1	0	0	0
5	3	4	2	0	0	0	1	0	0
0	5	1	6	6	0	0	0	1	0
$\backslash 5$	6	2	1	5	0	0	0	0	$_{1}/$

Experimentation

• If you want to see an article about estimating the average, variance, skewness, and kurtosis, for the number of occurrences of 1s, in row-echelon matrices of dimension $k \times 2k$ over GF(2), for k from 50 to 100, by simulating, 1000 times, look at the output file

http://www.math.rutgers.edu/~zeilberg/tokhniot/oCalabiWilf1.txt

• If you want to see an article about estimating the average, variance, skewness, and kurtosis, for the number of occurrences (as consecutive submatrix) of the matrix [[1, 0, 2], [1, 0, 2], [1, 0, 1]], in

row-echelon matrices of dimension $k \times 2k$ over GF(3) for k from 50 to 60 by simulating 1000 times, look at the output file

http://www.math.rutgers.edu/~zeilberg/tokhniot/oCalabiWilf1a.txt

• If you want to see an article about estimating the average, variance, skewness, and kurtosis of the number of 1s in $k \times 2k$ row-echelon matrices over GF(3) for $50 \le k \le 55$, by simulating 1000 times (and repeating each run three times, to compare notes), and comparing the estimated values to the (theoretically computed) exact values, look at

http://www.math.rutgers.edu/~zeilberg/tokhniot/oCalabiWilf2.txt

• If you want to see an article about estimating the average of the number of 1s in $k \times 3k$ row-echelon matrices over GF(2) for $100 \le k \le 110$, by simulating 1000 times (and repeating each run three times, to compare notes), and comparing the estimated average to the (theoretically computed) exact value, look at

http://www.math.rutgers.edu/~zeilberg/tokhniot/oCalabiWilf3.txt

• If you want to see an article on the average minimal weight of vectors in k-dimensional subspaces of $GF(2)^n$ for k from 1 to, 5, and n from 10 to 100, in increments of 10, look at the output file

http://www.math.rutgers.edu/~zeilberg/tokhniot/oCalabiWilf4.txt

For more output files see the web-page of this paper:

https://sites.math.rutgers.edu/~zeilberg/mamarim/mamarimhtml/calabi.html

Readers are welcome to experiment with our Maple package to their heart's content, and generate many more output files.

Conclusion: Rest in peace, my dear heroes, Herb and Gene, you both did great work on your own, and in collaboration with other people, in very different parts of mathematics, but this sole joint work is also great!

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Oct. 27, 2023.